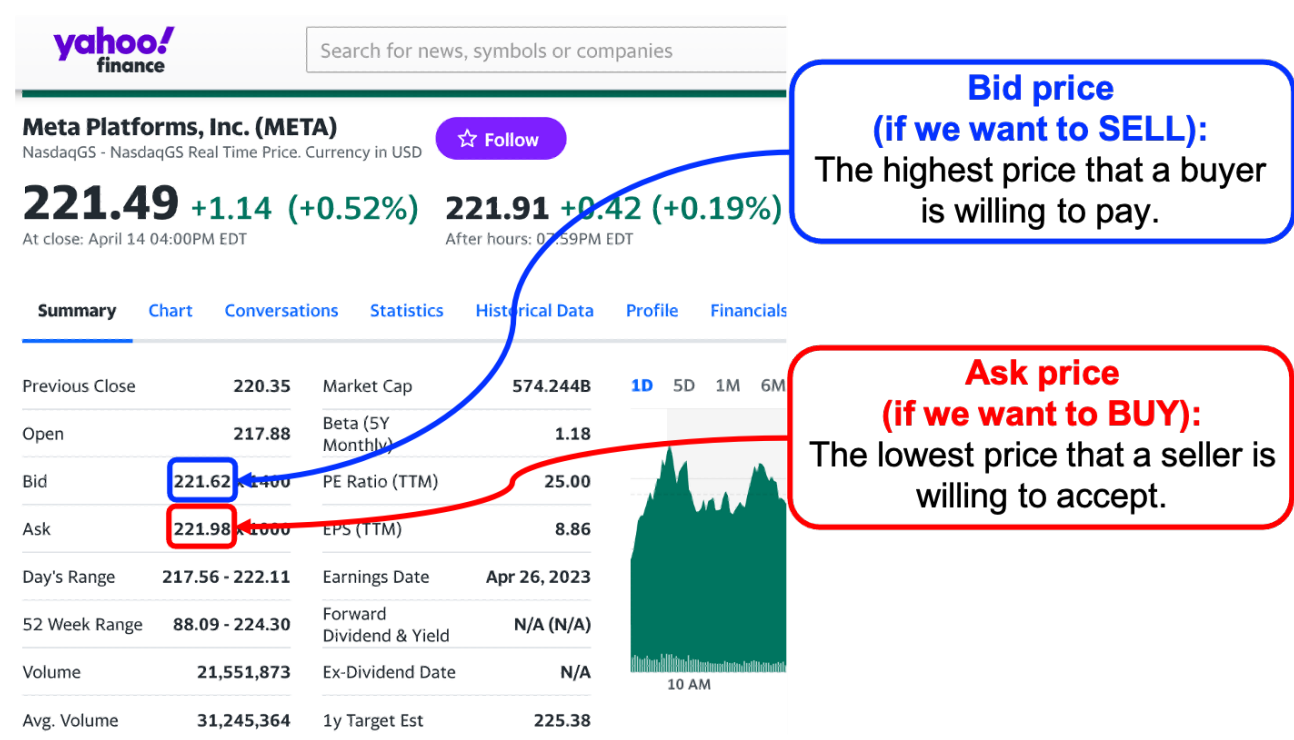


# No-Arbitrage Pricing with $\alpha$ -DS Mixtures in a Market with Bid-Ask Spreads

## Motivation: pricing in a market with frictions

- Classical no-arbitrage pricing theory assumes that the market is **competitive and frictionless**
- Prices can be expressed as **discounted expectations** with respect to an "artificial" probability measure  $Q$
- PROBLEM:** Markets show frictions, mostly in the form of **bid-ask spreads**
- AIM:** Replace  $Q$  with a non-additive measure so as to consider bid-ask spreads



## $\alpha$ -DS mixtures

Consider:

- $\Omega = \{1, \dots, n\}$  with  $n \geq 1$ , a finite set of states of the world
- $\mathcal{P}(\Omega)$ , power set of events
- $\mathbb{R}^\Omega$ , set of all random variables

### Definition ( $\alpha$ -DS mixture)

Let  $\alpha \in [0, 1]$ . A mapping  $\varphi_\alpha : \mathcal{P}(\Omega) \rightarrow [0, 1]$  is called an  $\alpha$ -DS mixture if there exists a belief function  $Bel : \mathcal{P}(\Omega) \rightarrow [0, 1]$  with dual plausibility function  $Pl$  such that, for all  $A \in \mathcal{P}(\Omega)$ ,

$$\varphi_\alpha(A) = \alpha Bel(A) + (1 - \alpha) Pl(A) = \alpha Bel(A) + (1 - \alpha)(1 - Bel(A^c)).$$

The belief function  $Bel$  is said to **represent** the  $\alpha$ -DS mixture  $\varphi_\alpha$ .

We further distinguish the subclasses of **additive and consonant**  $\alpha$ -DS mixtures.

### Proposition (unique representation)

Let  $\alpha \in [0, 1]$  with  $\alpha \neq \frac{1}{2}$ , and  $\varphi_\alpha : \mathcal{P}(\Omega) \rightarrow [0, 1]$  be an  $\alpha$ -DS mixture. Let  $Bel, Bel'$  be belief functions on  $\mathcal{P}(\Omega)$ . If both  $Bel$  and  $Bel'$  represent  $\varphi_\alpha$ , then  $Bel = Bel'$ .

## Properties of $\alpha$ -DS mixtures

### Proposition (properties of a $\varphi_\alpha$ )

Let  $\alpha \in [0, 1]$ . An  $\alpha$ -DS mixture  $\varphi_\alpha : \mathcal{P}(\Omega) \rightarrow [0, 1]$  satisfies the following properties:

- $\varphi_\alpha(\emptyset) = 0$  and  $\varphi_\alpha(\Omega) = 1$ ;
- $\varphi_\alpha(A) \leq \varphi_\alpha(B)$ , when  $A \subseteq B$  and  $A, B \in \mathcal{P}(\Omega)$ ;
- $\varphi_\alpha$  is self-dual if and only if it is additive or  $\alpha = \frac{1}{2}$ ;
- $\varphi_\alpha$  is sub-additive if it is additive or  $\alpha \in [0, \frac{1}{2}]$ .

For every  $\alpha \in [0, 1]$ , the class  $\mathcal{M}_\alpha$  of all  $\alpha$ -DS mixtures on  $\mathcal{P}(\Omega)$  is **convex** and contains the class  $\mathcal{P}$  of all probability measures on  $\mathcal{P}(\Omega)$ .

## $\alpha$ -DS mixture Choquet expectation

Every  $\varphi_\alpha$  uniquely extends to a functional  $C_{\varphi_\alpha} : \mathbb{R}^\Omega \rightarrow \mathbb{R}$  by setting, for every  $X \in \mathbb{R}^\Omega$ ,

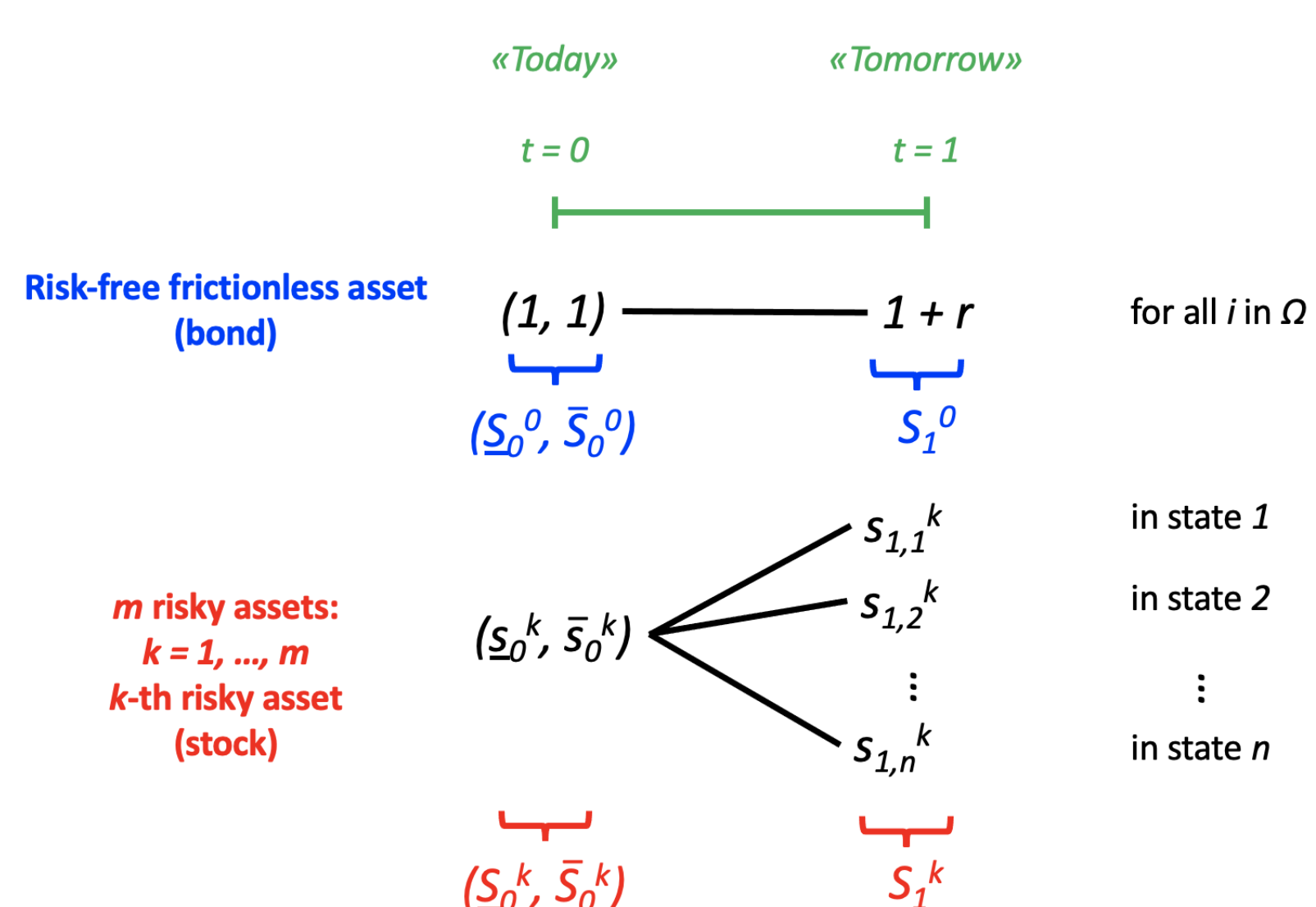
$$C_{\varphi_\alpha}[X] = \int X d\varphi_\alpha$$

**Hurwicz-like representation:**  $C_{\varphi_\alpha}[X] = \alpha \min_{P \in C_{Bel}} \mathbb{E}_P[X] + (1 - \alpha) \max_{P \in C_{Bel}} \mathbb{E}_P[X]$  where  $C_{Bel}$  is the core of  $Bel$

**Möbius-like representation:**  $C_{\varphi_\alpha}[X] = \sum_{B \in \mathcal{U}} [X]^\alpha(B) \mu(B)$  where  $\mu$  is the Möbius inverse of  $Bel$  and

$$\mathcal{U} = \mathcal{P}(\Omega) \setminus \{\emptyset\} \quad \text{and} \quad [X]^\alpha : \mathcal{U} \rightarrow \mathbb{R} \quad \text{with} \quad [X]^\alpha(B) = \alpha \min_{i \in B} X(i) + (1 - \alpha) \max_{i \in B} X(i)$$

## One-period market with bid-ask spreads



$\alpha$ -Mixture prices at time  $t = 0$ :

$$M_0 = (S_0^0, S_0^1, \dots, S_0^m)^T$$

$$S_0^k = \alpha S_0^k + (1 - \alpha) \bar{S}_0^k$$

Payoffs at time  $t = 1$ :

$$M_1 = (S_1^0, S_1^1, \dots, S_1^m)^T$$

## No-arbitrage pricing under $\alpha$ -PRU

Given a portfolio  $\lambda = (\lambda_0, \lambda_1, \dots, \lambda_m)^T \in \mathbb{R}^{m+1}$  we define:

$$\text{Price at time } t = 0: V_0^\lambda = \lambda_0 + \sum_{k=1}^m \lambda_k S_0^k$$

$$\text{Payoff under } \alpha\text{-PRU at time } t = 1: V_1^\lambda = \lambda_0(1+r) + \sum_{k=1}^m \lambda_k [S_1^k]^\alpha$$

### $\alpha$ -PRU principle at time $t = 1$

**PRU (Partially Resolving Uncertainty):** An agent may only acquire that  $B \neq \emptyset$  occurs, without knowing which is the true  $i \in B$

**$\alpha$ -pessimism:** An agent always considers the  $\alpha$ -mixture between the minimum and the maximum of random payoffs on every  $B \neq \emptyset$

### Theorem (First FTAP under $\alpha$ -PRU)

Let  $\alpha \in [0, 1]$ . The following conditions are equivalent:

- there exists an  $\alpha$ -DS mixture  $\widehat{\varphi}_\alpha$  represented by a belief function strictly positive on  $\mathcal{U}$  and such that  $\frac{C_{\widehat{\varphi}_\alpha}[S_1^k]}{1+r} = S_0^k$ , for  $k = 1, \dots, m$ ;
- for every  $\lambda \in \mathbb{R}^{m+1}$  none of the following conditions holds:
  - $V_1^\lambda(\{i\}) = 0$ , for  $i = 1, \dots, n$ ,  $V_1^\lambda(B) \geq 0$ , for all  $B \in \mathcal{U} \setminus \{\{i\} : i \in \Omega\}$  and  $V_0^\lambda < 0$ ;
  - $V_1^\lambda(\{i\}) \geq 0$ , for  $i = 1, \dots, n$ , not all 0,  $V_1^\lambda(B) \geq 0$ , for all  $B \in \mathcal{U} \setminus \{\{i\} : i \in \Omega\}$ , and  $V_0^\lambda \leq 0$ .

### Theorem (Second FTAP under $\alpha$ -PRU)

Let  $\alpha \in [0, 1]$ . If the market satisfies condition (ii) of the First FTAP under  $\alpha$ -PRU and is  $\alpha$ -PRU complete, i.e., for  $\mathcal{U} = \{B_1, \dots, B_{2^n-1}\}$ , it is  $m \geq 2^n - 1$  and  $S_1^k = 1_{B_k}$ , for  $k = 1, \dots, 2^n - 1$ , then the  $\alpha$ -DS mixture  $\widehat{\varphi}_\alpha$  in condition (i) of the First FTAP under  $\alpha$ -PRU is unique.

### $\alpha$ -DS mixture no-arbitrage price of a payoff $X_1 \in \mathbb{R}^\Omega$

$$X_0 = (1+r)^{-1} C_{\widehat{\varphi}_\alpha}[X_1] = (1+r)^{-1} \left( \alpha \min_{Q \in C_{\widehat{Bel}}} \mathbb{E}_Q[X_1] + (1-\alpha) \max_{Q \in C_{\widehat{Bel}}} \mathbb{E}_Q[X_1] \right)$$

## META stock market data with bid-ask spreads

Consider a single risky asset:

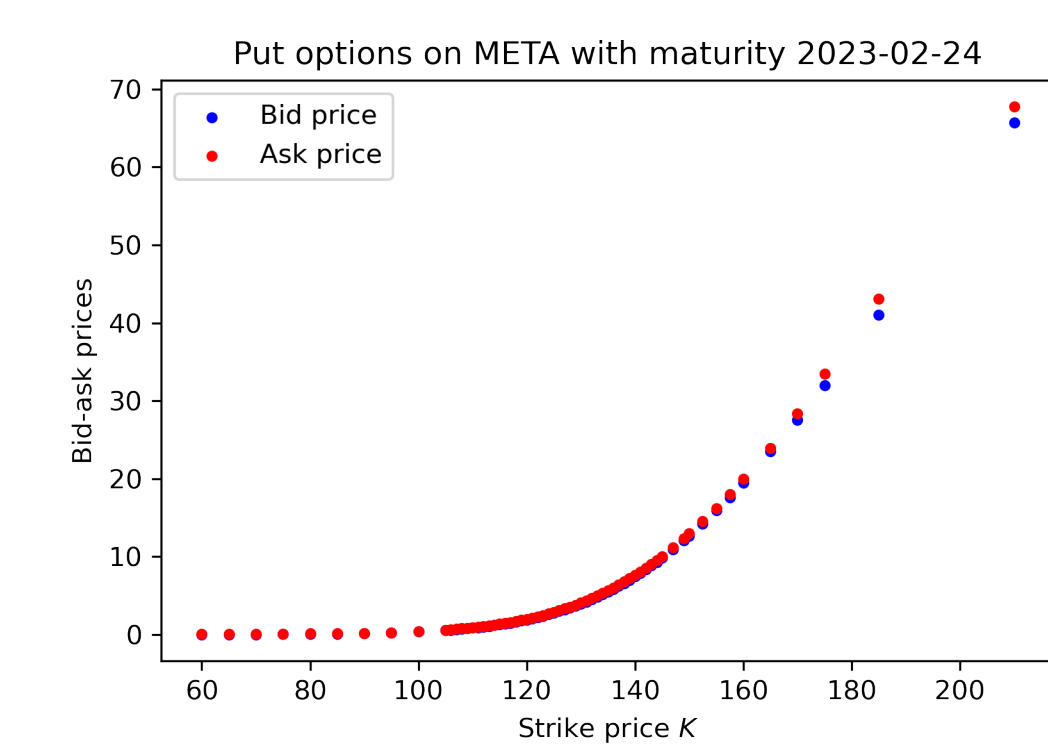
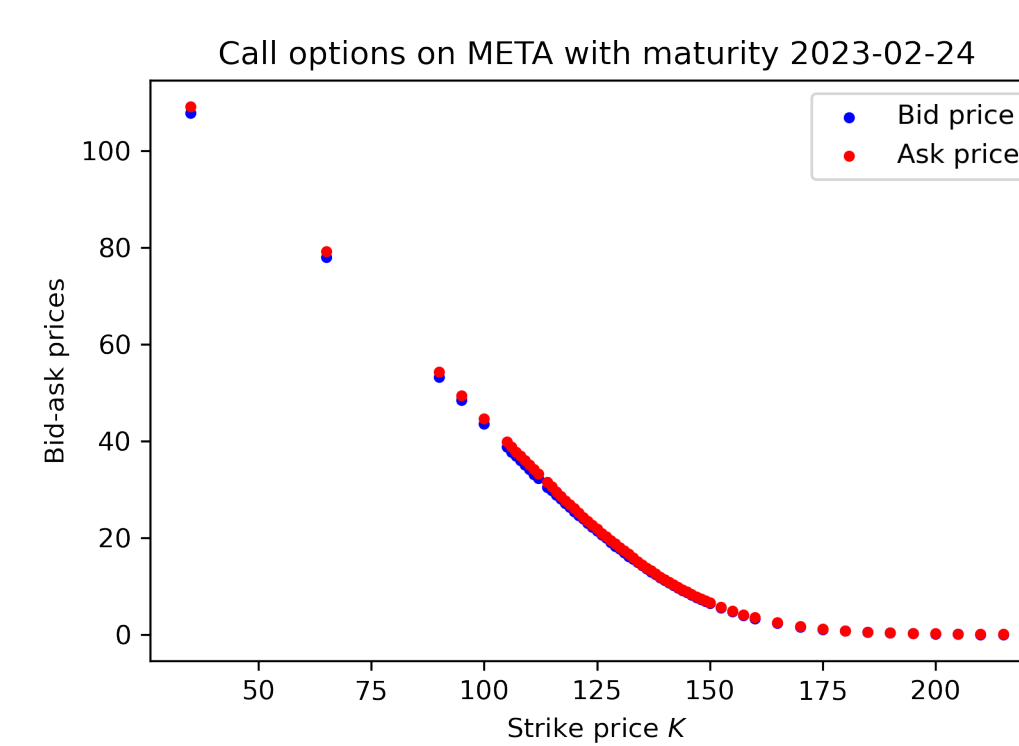
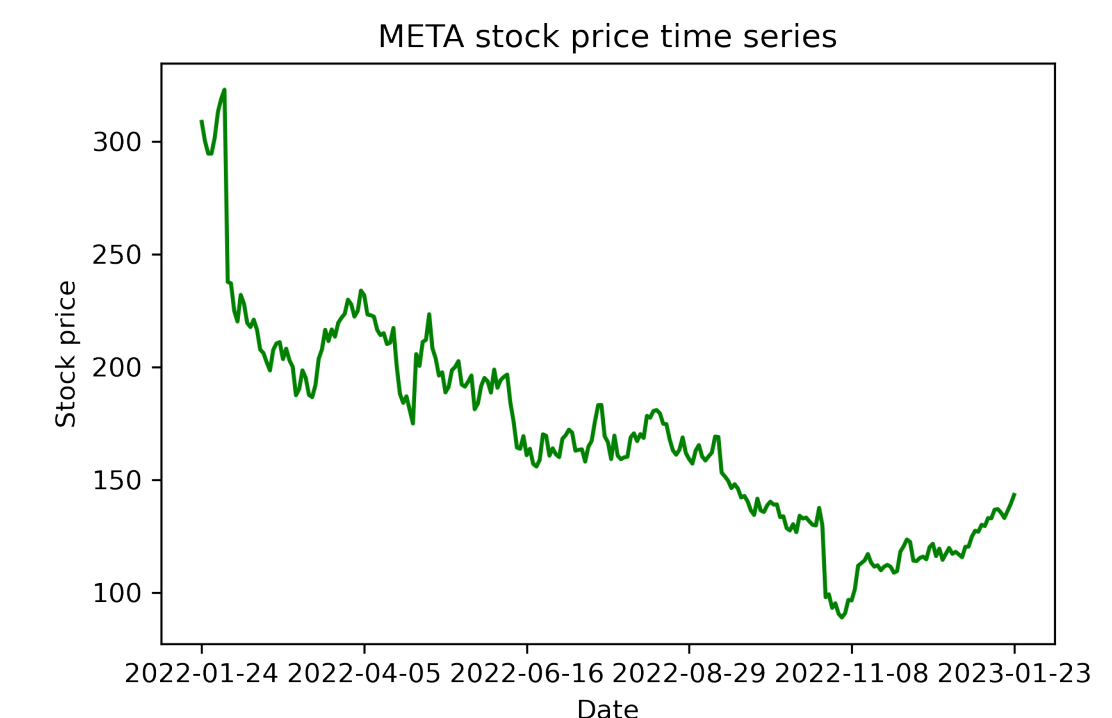
- $t = 0$  identified with 2023-01-23
- $t = 1$  identified with 2023-02-24
- US T-Bill with  $1+r = (1.0469)^{\frac{32}{365}}$
- Last one year of META closing prices:

$$S_1^1 \text{ ranging in } [112.4, 159.2, 206.0, 252.8, 299.6]$$

- Bid-ask prices at time  $t = 0$  of call and put options on META with maturity  $t = 1$ , strike prices in  $\mathcal{K}_{call}$  and  $\mathcal{K}_{put}$ , and payoffs

$$C^K = \max\{S_1^1 - K\}$$

$$P^K = \max\{K - S_1^1\}$$



## Tuning of $\alpha$ : a measure of market pessimism

For a fixed  $\alpha \in [0, 1]$ , compute the  $\alpha$ -mixture prices  $C_0^{K,\alpha} = \alpha C_0^K + (1-\alpha) \bar{C}_0^K$  and  $P_0^{K,\alpha} = \alpha P_0^K + (1-\alpha) \bar{P}_0^K$ :

$$\text{minimize } E(\widehat{\varphi}_\alpha) = \sum_{K \in \mathcal{K}_{call}} \left( C_0^{K,\alpha} - \frac{C_0^K}{1+r} \right)^2 + \sum_{K \in \mathcal{K}_{put}} \left( P_0^{K,\alpha} - \frac{P_0^K}{1+r} \right)^2$$

subject to:

$$\begin{cases} \widehat{\varphi}_\alpha \in \mathcal{M}_\alpha, \\ \widehat{\varphi}_\alpha \text{ is represented by } \widehat{Bel}, \\ \widehat{Bel}(\{i\}) \geq \epsilon, \text{ for all } i \in \Omega, \text{ with } \epsilon = 0.0001 \end{cases}$$

