

A Nonstandard Approach To Stochastic Processes Under Probability Bounding

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What is internal set theory?

- ▶ new predicate '**standard**' that applies to objects (sets, functions, ...)
- ▶ three new axioms added to ZFC to govern use of this predicate [4]

Intuition?

- ▶ internal formulas: do not use 'standard' "0 + 1 equals 1"
- ▶ external formulas: do use 'standard' "0 is a standard natural number"
- ▶ for an object to be standard, intuitively, we mean [2, §1.1.1, p. 2]:
'at any stage within the mathematical discourse, [...] uniquely defined [using an explicitly written internal formula]'

Why is it useful?

- ▶ has the notion of an **infinitesimal** so the theory formalizes intuition that goes back to Newton, Leibniz, Cauchy, etc. [1]
- ▶ many objects have an infinitely close standard object, called its **shadow**: allows us to move between standard functions with infinite domain and non-standard functions with finite domain
very useful to study stochastic processes in continuous time!

Difficulties

- ▶ for historical reasons, most mathematicians are unfamiliar with it
- ▶ application of new axioms needs care: requires some retraining in logic
- ▶ illegal set formation: cannot form sets with external formulas

What is a stochastic process?

Notation:

- ▶ standard finite state space \mathcal{X}
- ▶ index set $T_0 \subseteq \mathbb{R}_+$
- ▶ possibility space Ω
- ▶ function $\xi: T_0 \times \Omega \rightarrow \mathcal{X}$ given outcome ω at time t , the state is $\xi(t, \omega)$
- ▶ $T \subseteq T_0$ any subset of T_0 (e.g. finite discretization)
- ▶ $\mathcal{A}(T)$ algebra generated by events of the form $\{\xi(t) = x\}$
- ▶ $\mathcal{L}(T)$ linear space spanned by $\mathcal{A}(T)$
- ▶ $\mathcal{K}(T) := \mathcal{L}(T) \times (\mathcal{A}(T) \setminus \{\emptyset\})$

Definition

A (stochastic) process on T is a coherent lower prevision \mathbb{E} defined on $\mathcal{K}(T)$.

Nearby elementary processes

Theorem

Assume $T \subseteq T_0$ and T contains all standard elements of T_0 .

Let $\mathbb{E}: \mathcal{K}(T) \rightarrow \mathbb{R}$ be any process.

Then there is a unique standard process $\mathbb{E}_0: \mathcal{K}(T_0) \rightarrow \mathbb{R}$,

called the **shadow** of \mathbb{E} , satisfying

$$\forall (f, A) \in \mathcal{K}(T_0): \mathbb{E}(f | A) \simeq \mathbb{E}_0(f | A) \quad (1)$$

Definition

A standard process \mathbb{E}_0 on T_0 is said to be *nearby* an elementary process \mathbb{E} on T if $T \subseteq T_0$, T contains all standard elements of T_0 , and \mathbb{E}_0 is the shadow of \mathbb{E} .

Imprecise Markov chains: Notation

- ▶ T is a finite subset of T_0 containing all standard elements of T_0
- ▶ $T' := T \setminus \{\max T\}$
- ▶ if $t \in T'$ then $t + dt$ denotes the successor of t in T , i.e.
$$dt := \min\{t' \in T: t' > t\} - t \quad (2)$$
- ▶ for any function ϕ on T , $\phi(0 : t)$ denotes the restriction of ϕ to $[0, t] \cap T$

Imprecise Markov chains: Definition

- ▶ \mathbb{I} : coherent lower prevision on $\mathbb{R}^{\mathcal{X}}$
- ▶ $\mathbb{I}_t(\cdot)(x)$: coherent lower prevision on $\mathbb{R}^{\mathcal{X}}$ (for each $t \in T'$ and $x \in \mathcal{X}$)

Definition

A precise elementary process \mathbb{E} on T is **compatible** with (\mathbb{I}, \mathbb{I}) if for all paths $x: T' \rightarrow \mathcal{X}$, all $t \in T'$, and all $f \in \mathbb{R}^{\mathcal{X}}$,

$$\mathbb{E}(f(\xi(0))) \geq \mathbb{I}(f) \quad (3)$$

$$\mathbb{E}(f(\xi(t + dt)) | \xi(0 : t) = x(0 : t)) \geq \mathbb{I}_t(f)(x(t)) \quad (4)$$

If $\underline{\mathbb{E}}$ denotes the lower envelope of all these compatible precise elementary processes, then $\underline{\mathbb{E}}$ is called the **elementary imprecise Markov chain** induced by (\mathbb{I}, \mathbb{I}) . Its shadow is called the **imprecise Markov chain** induced by (\mathbb{I}, \mathbb{I}) .

Continuous time imprecise Markov chains

- ▶ fix a standard **lower rate operator** $\underline{\mathbb{Q}}: \mathbb{R}^{\mathcal{X}} \rightarrow \mathbb{R}^{\mathcal{X}}$ and define

$$\mathbb{T}_t := I + dt\underline{\mathbb{Q}} \quad (5)$$

- ▶ the shadow \mathbb{E}_0 of this elementary imprecise Markov chain behaves just like the usual continuous time imprecise Markov chains

Theorem

For all $x \in \mathcal{X}$, $f \in \mathbb{R}^{\mathcal{X}}$, and $t \geq 0$, $\mathbb{E}_0(f(\xi(t)) | \xi(0) = x) = e^{t\underline{\mathbb{Q}}}(f)(x)$.

Sketch of proof

Assume $t > 0$. By transfer, only need to establish the equality for standard t . It suffices to show that

$$\left\| \prod_{s < t} (I + ds\underline{\mathbb{Q}}) - e^{t\underline{\mathbb{Q}}} \right\| \simeq 0 \quad (6)$$

Indeed, with $\delta := \max_{s < t} ds$,

$$\left\| \prod_{s < t} (I + ds\underline{\mathbb{Q}}) - e^{t\underline{\mathbb{Q}}} \right\| = \left\| \prod_{s < t} (I + ds\underline{\mathbb{Q}}) - \prod_{s < t} e^{ds\underline{\mathbb{Q}}} \right\| \quad (7)$$

$$\leq \sum_{s < t} \left\| I + ds\underline{\mathbb{Q}} - e^{ds\underline{\mathbb{Q}}} \right\| \lesssim \sum_{s < t} (ds)^2 \|\underline{\mathbb{Q}}\|^2 \quad (8)$$

$$\leq \sum_{s < t} (ds)\delta \|\underline{\mathbb{Q}}\|^2 = t\delta \|\underline{\mathbb{Q}}\|^2 \simeq 0 \quad (9)$$

since t and $\|\underline{\mathbb{Q}}\|$ are limited, and $\delta \simeq 0$.

(compare with [3])

[1] Alexandre Borovik and Mikhail G. Katz, *Who gave you the Cauchy–Weierstrass tale? The dual history of rigorous calculus*, Foundations of Science 17 (2012), no. 3, 245–276.

[2] Francine Diener and Marc Diener (eds.), *Nonstandard analysis in practice*, Springer, 1995.

[3] Thomas Krak, Jasper De Bock, and Arno Siebes, *Imprecise continuous-time Markov chains*, International Journal of Approximate Reasoning 88 (2017), 452–528.

[4] Edward Nelson, *Radically elementary probability theory*, Annals of Mathematical Studies, Princeton University Press, New Jersey, 1987.